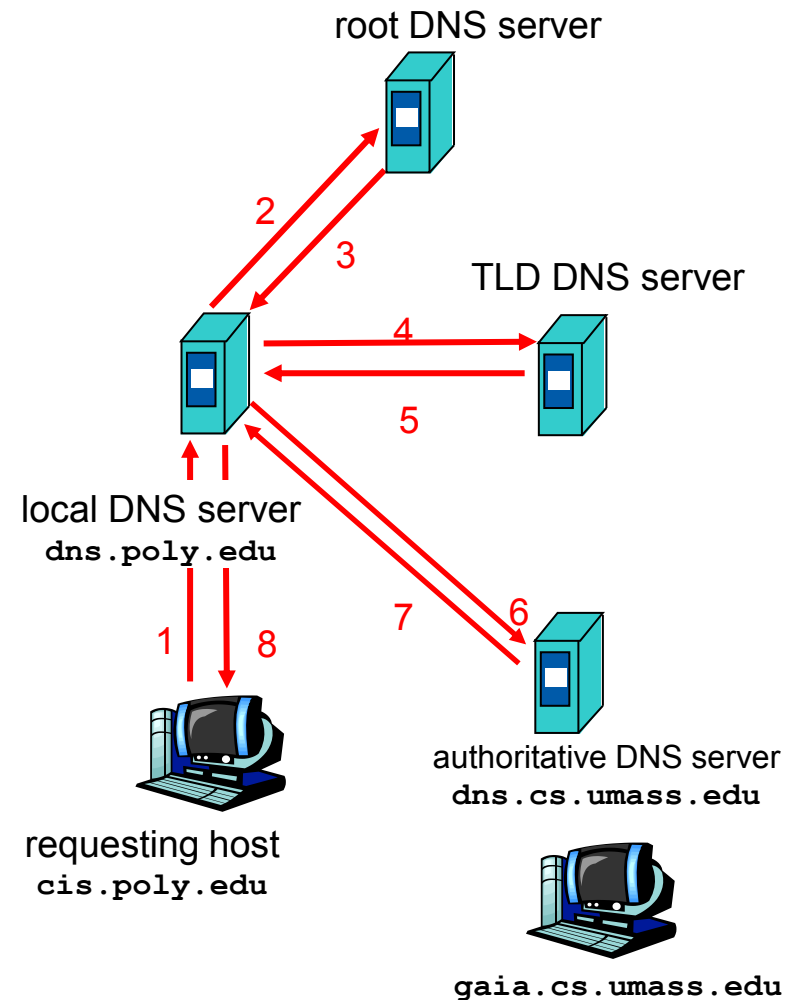


DNS – Example of Iterative Queries



- Host at `cis.poly.edu` wants IP address for **`gaia.cs.umass.edu`**



DNS – Recursive Queries

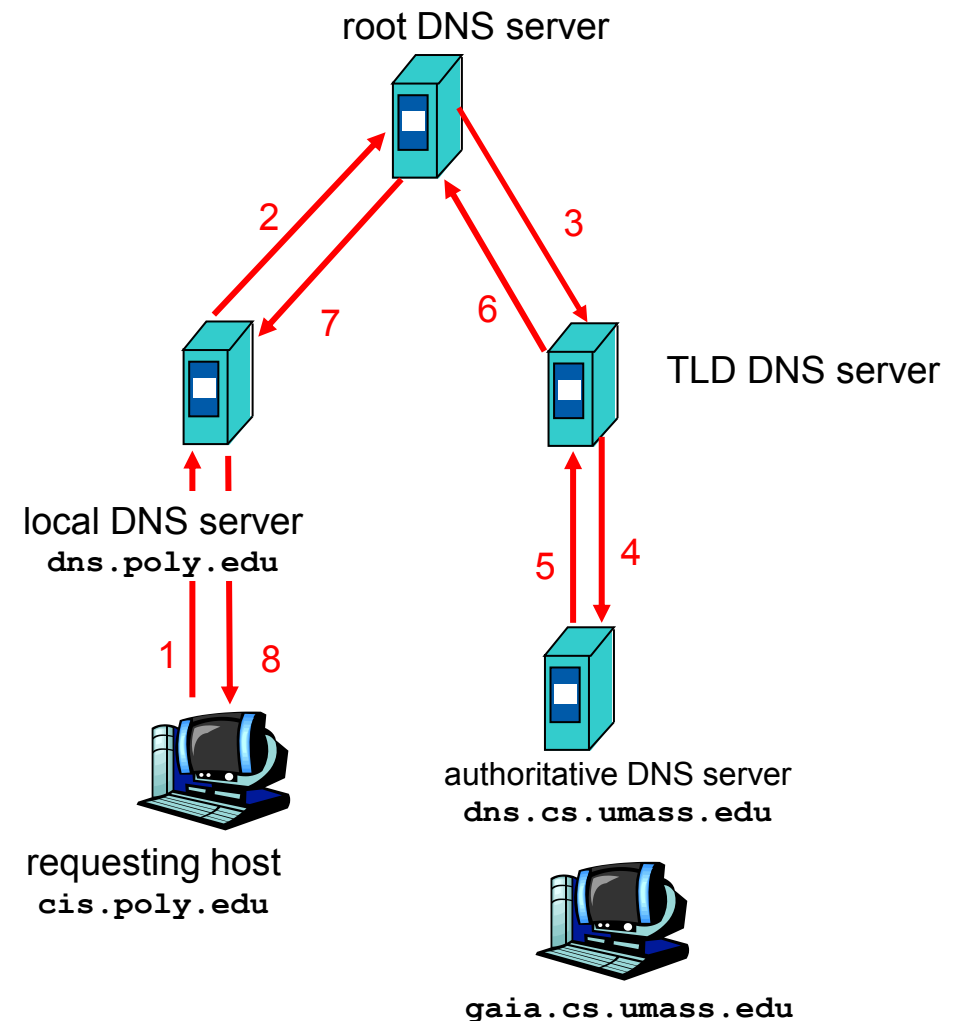


Recursive query:

- Puts burden of name resolution on contacted name server
- Heavy load?

Iterated query:

- Contacted server replies with name of server to contact
- “I don’t know this name, but ask this server”



DNS: Caching and Updating Records



- Once (any) name server learns mapping, it *caches* mapping
 - Cache entries timeout (disappear) after some time
 - TLD servers typically cached in local name servers
 - Thus root name servers not often visited

- Update/notify mechanisms under design by IETF
 - RFC 2136
 - <http://www.ietf.org/html.charters/dnsind-charter.html>



Inserting Records Into DNS

- Example: just created startup “Network Utopia”
- Register name networkutopia.com at a registrar (e.g., Network Solutions)
 - Need to provide registrar with names and IP addresses of your authoritative name server (**primary** and **secondary**)
 - Registrar inserts two RRs into the com TLD server:

(networkutopia.com, dns1.networkutopia.com, NS)
(dns1.networkutopia.com, 212.212.212.1, A)

- Put in authoritative server Type A record for www.networkutopia.com and Type MX record for networkutopia.com



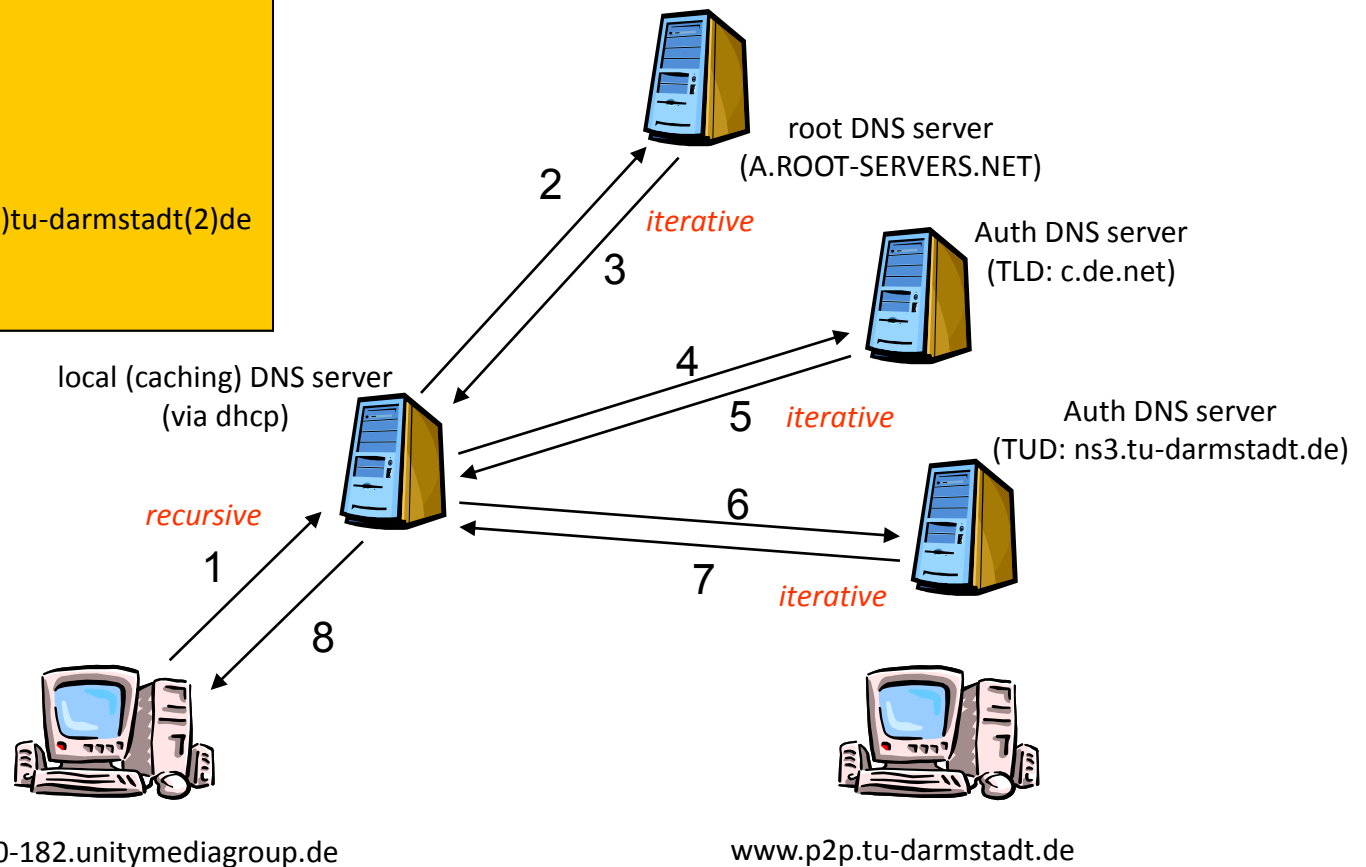
DNS – Recursive and Iterative Queries

DNS HEADER (send)

- Identifier: 0x3116
- Flags: 0x00 (Q)
- Opcode: 0 (Standard query)
- Return code: 0 (No error)
- Number questions: 1
- Number answer RR: 0
- Number authority RR: 0
- Number additional RR: 0

QUESTIONS (send)

- Queryname: (3)www(3)p2p(12)tu-darmstadt(2)de
- Type: 1 (A)
- Class: 1 (Internet)





A Quick Example...

```
strufe@eris:~$ dnstracer -v www.p2p.tu-darmstadt.de
```

```
Tracing to informatik.tu-darmstadt.de[a] via 130.83.163.141, maximum of 3 retries
```

```
130.83.163.141 (130.83.163.141) IP HEADER
```

```
-Destination address: 130.83.163.141
```

```
-DNS HEADER (send)
```

```
-- Identifier:      0x3116
```

```
-- Flags:          0x00 (Q )
```

```
-- Opcode:         0 (Standard query)
```

```
-- Return code:    0 (No error)
```

```
-- Number questions: 1
```

```
-- Number answer RR: 0
```

```
-- Number authority RR: 0
```

```
-- Number additional RR: 0
```

```
-QUESTIONS (send)
```

```
-- Queryname:      (3)www(3)p2p(12)tu-darmstadt(2)de
```

```
-- Type:           1 (A)
```

```
-- Class:          1 (Internet)
```

```
-DNS HEADER (recv)
```

```
-- Identifier:      0x3116
```

```
-- Flags:          0x8080 (R RA )
```

```
-- Opcode:         0 (Standard query)
```

```
-- Return code:    0 (No error)
```

```
-- Number questions: 1
```

```
-- Number answer RR: 2
```

```
-- Number authority RR: 0
```

```
-- Number additional RR: 0
```

```
-.....
```

```
QUESTIONS (recv)
```

```
- Queryname:      (3)www(3)p2p(12)tu-darmstadt(2)de
```

```
- Type:           1 (A)
```

```
- Class:          1 (Internet)
```

```
ANSWER RR
```

```
- Domainname:     (6)charon(7)dekanat(10)informatik(12)tu-darmstadt(2)de
```

```
- Type:           1 (A)
```

```
- Class:          1 (Internet)
```

```
- TTL:            1592 (26m32s)
```

```
- Resource length: 4
```

```
- Resource data:   130.83.162.6
```

```
ANSWER RR
```

```
- Domainname:     (3)www(3)p2p(12)tu-darmstadt(2)de
```

```
- Type:           5 (CNAME)
```

```
- Class:          1 (Internet)
```

```
- TTL:            49817 (13h50m17s)
```

```
- Resource length: 28
```

```
- Resource data:   (6)charon(7)dekanat(10)informatik(12)tu-darmstadt(2)de
```

```
Got answer [received type is cname]
```



So where is the Info?

```
strufe@eris:~$ dnstracer -v -qns tu-darmstadt.de
```

```
Tracing to tu-darmstadt.de[ns] via 130.83.163.130
```

```
130.83.163.130 (130.83.163.130) IP HEADER
```

```
- Destination address: 130.83.163.130
```

```
DNS HEADER (send)
```

```
- Identifier: 0x4C45
```

```
- Flags: 0x00 (Q )
```

```
- Opcode: 0 (Standard query)
```

```
- Return code: 0 (No error)
```

```
- Number questions: 1
```

```
- Number answer RR: 0
```

```
- Number authority RR: 0
```

```
- Number additional RR: 0
```

```
QUESTIONS (send)
```

```
- Queryname: (12)tu-darmstadt(2)de
```

```
- Type: 2 (NS)
```

```
- Class: 1 (Internet)
```

```
DNS HEADER (recv)
```

```
- Identifier: 0x4C45
```

```
- Flags: 0x8080 (R RA )
```

```
- Opcode: 0 (Standard query)
```

```
- Return code: 0 (No error)
```

```
- Number questions: 1
```

```
- Number answer RR: 5
```

```
- Number authority RR: 0
```

```
- Number additional RR: 9
```

```
.....
```

```
QUESTIONS (recv)
```

```
- Queryname: (12)tu-darmstadt(2)de
```

```
- Type: 2 (NS)
```

```
- Class: 1 (Internet)
```

```
ANSWER RR
```

```
- Domainname: (12)tu-darmstadt(2)de
```

```
- Type: 2 (NS)
```

```
- Class: 1 (Internet)
```

```
- TTL: 70523 (19h35m23s)
```

```
- Resource length: 6
```

```
- Resource data: (3)ns1(3)hrz(12)tu-darmstadt(2)de
```

```
ANSWER RR
```

```
- Domainname: (12)tu-darmstadt(2)de
```

```
- Type: 2 (NS)
```

```
- Class: 1 (Internet)
```

```
- TTL: 70523 (19h35m23s)
```

```
- Resource length: 5
```

```
- Resource data: (2)ns(6)man-da(2)de
```

```
ANSWER RR
```

```
- Domainname: (12)tu-darmstadt(2)de
```

```
- Type: 2 (NS)
```

```
- Class: 1 (Internet)
```

```
- TTL: 70523 (19h35m23s)
```

```
- Resource length: 6
```

```
- Resource data: (3)ns2(3)hrz(12)tu-darmstadt(2)de
```

```
.....
```

Answer ctd...



```
.....
ADDITIONAL RR
- Domainname:      (3)ns1(3)hrz(12)tu-darmstadt(2)de
- Type:            1 (A)
- Class:           1 (Internet)
- TTL:             17335 (4h48m55s)
- Resource length:  4
- Resource data:    130.83.22.63
ADDITIONAL RR
- Domainname:      (2)ns(6)man-da(2)de
- Type:            28 (unknown)
- Class:           1 (Internet)
- TTL:             38386 (10h39m46s)
- Resource length:  16
- Resource data:    2001:41b8:0000:0001:0000:0000:0000:0053
ADDITIONAL RR
- Domainname:      (2)ns(6)man-da(2)de
- Type:            1 (A)
- Class:           1 (Internet)
- TTL:             38386 (10h39m46s)
- Resource length:  4
- Resource data:    82.195.66.249
ADDITIONAL RR
- Domainname:      (3)ns2(3)hrz(12)tu-darmstadt(2)de
- Type:            28 (unknown)
- Class:           1 (Internet)
- TTL:             17335 (4h48m55s)
- Resource length:  16
- Resource data:    2001:41b8:083f:0022:0000:0000:0000:0063
.....
```

```
.....
ADDITIONAL RR
- Domainname:      (3)ns2(3)hrz(12)tu-darmstadt(2)de
- Type:            1 (A)
- Class:           1 (Internet)
- TTL:             17335 (4h48m55s)
- Resource length:  4
- Resource data:    130.83.22.60
ADDITIONAL RR
- Domainname:      (3)ns2(6)man-da(2)de
- Type:            1 (A)
- Class:           1 (Internet)
- TTL:             38386 (10h39m46s)
- Resource length:  4
- Resource data:    217.198.242.225
ADDITIONAL RR
- Domainname:      (3)ns3(3)hrz(12)tu-darmstadt(2)de
- Type:            28 (unknown)
- Class:           1 (Internet)
- TTL:             17335 (4h48m55s)
- Resource length:  16
- Resource data:    2001:41b8:083f:0056:0000:0000:0000:0060
ADDITIONAL RR
- Domainname:      (3)ns3(3)hrz(12)tu-darmstadt(2)de
- Type:            1 (A)
- Class:           1 (Internet)
- TTL:             17335 (4h48m55s)
- Resource length:  4
- Resource data:    130.83.56.60
Got answer
```

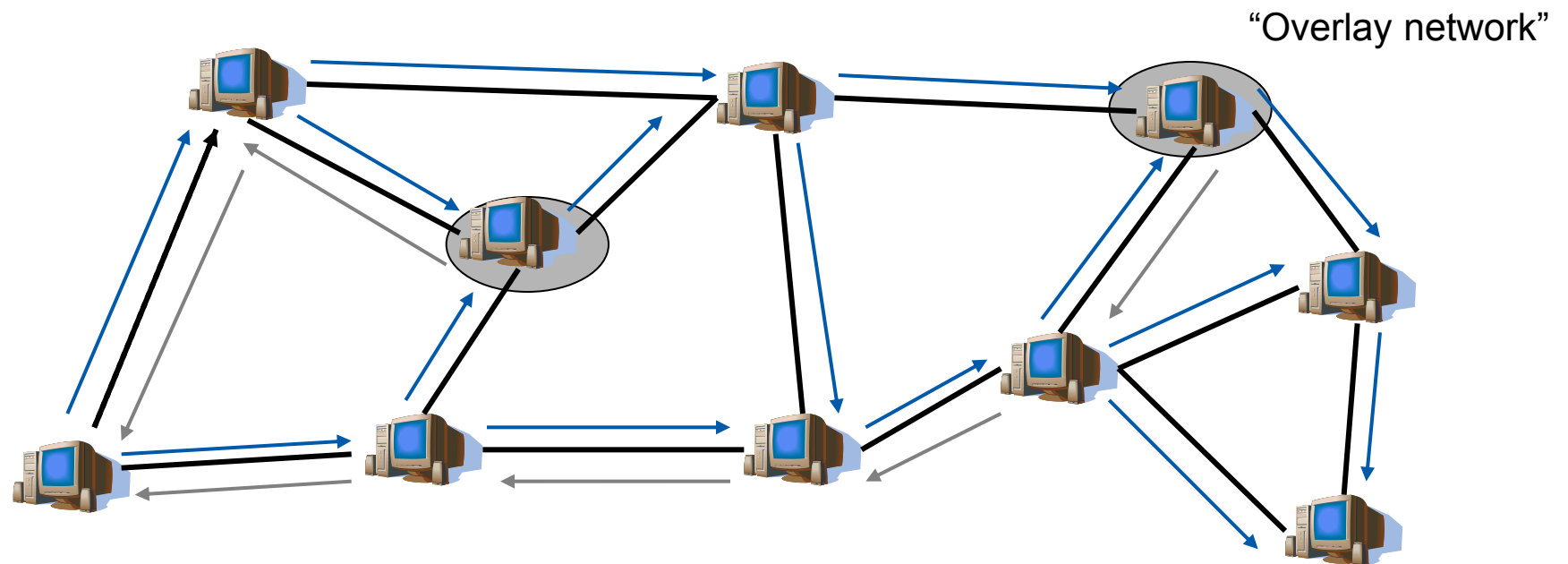



DNS – Lessons Learned

1. Structure name space (divide et impera)
2. Simple „routing“ b/c of structured (hierarchical) namespace
3. Store information at multiple locations
4. Maintain multiple connections
5. Be redundant! (Replicate...)
 - primary and secondary server, multiple TLD servers
6. Delegation using iterative or recursive forwarding
(Btw: what are the pros and cons of each?)

What does this „routing“ mean anyways!?

Back to P2P!





P2P in a Nutshell

- Properties of (pure) P2P: ***„All peers are equal“***
 - no dedicated service, no central entity
 - no a-priori knowledge / structure / hierarchy
 - highly dynamic behavior of nodes→ Flat system architecture, flat namespace, unreliable service providers

- Main primary problems of P2P:
 - Staying connected
 - Resource lookup (name resolution, location of replica, especially selecting a good next hop for the delegation → routing)
 - Can't trust anyone



Peer-to-Peer (a Definition)

- **Communication model:** asynchronous (request-response)
- **Role model:** a single role (?)
 - symmetric behavior, all peers in general (can) do the same
 - **BUT:** considering an interaction there is one requesting and n responding peers.
- **Organisational model:** completely unstructured („it's a mess!")
 - Other than bootstrapping no knowledge whatsoever about the context, no knowledge about the structure
- No **Identifiers**, only names
- We can introduce identifiers based on distributed algorithms (hashes)
- We can introduce structure using distributed algorithms (supernodes, etc.)
- *A P2P overlay on the Internet is a subset of links of a clique graph*

The P2P Environment



...all this in order to do:

File sharing, content distribution (BitTorrent/iptv), session initiation/chat/voip (skype, jabber), malware distribution/spam (botnets),...

- Standard Solutions (p2p – the executive summary)
 - *Connectivity*: select enough fall-back „servers“
 - *Name resolution*: unstructured P2P (flooding) or external search engine
 - *Resource location*: registry and lookup in structured P2P (DHT!) (or the above...)
- Closely related fields
 - Delay Tolerant Networks (Ad-hoc-, opportunistic-, pocket-switched-, vehicular-, <you-name-it> networks)
 - Wireless Sensor Networks
 - Epidemic-, Content-/ Context-based routing



Properties of P2P Systems

- P2P systems typically have the following properties:
 1. **Unreliable, uncoordinated, unmanaged**
 - No central authority, peers are completely independent
 - Increases flexibility of individual peers, but makes the overall system (possibly) unreliable
 2. **Resilient to attacks, heterogeneous**
 - Large number of peers in the system, hard to bring it down?
 - Heterogeneous peers make viruses and worms harder to write?
 3. **Large collection of resources**
 - Voluntary participation, global reach
 - Millions of simultaneous users



History of P2P

- What was the first P2P system and when?
- Answer: ARPANet 1969
- Later: USENET, 1979 (also FidoNet 1984, other BBSs)
 - Current Internet routing (BGP) is P2P
- The term P2P was coined by Napster in 1999
- Napster was a huge hit, brought P2P to general attention
- Illegal sharing of copyrighted material by users was the main driver behind Napster's success and the reason for its downfall
- Other systems followed Napster quickly, based on other design choices
- Research community followed suit quickly
 - Many deployed systems proprietary, hard to examine well...



Current State of P2P

- Where are we now?
- P2P networks going strong, all over the world
 - Many networks highly popular and widely used
 - Different networks in different countries
- P2P networks currently mostly used for illegal sharing of copyrighted material
 - Music, videos, software, ...
 - Note: Can be used for legal sharing too (see BitTorrent)
- Other applications starting to emerge (see below)
- Content providers not so happy
 - Sue companies making P2P software (e.g., Napster), sue software developers (Winny), sue users sharing material
 - But also providing alternate means: iTunes & friends



New P2P Systems

- File sharing was first P2P application
- Other applications are coming to light
- BitTorrent more content distribution than file sharing
- P2P extending beyond file sharing: [Skype](#)
 - We will look at Skype closer in Chapter 2
- Skype is a P2P telephone “system”
 - Can call other computers, or normal phones
- Skype is based on the KaZaA network (see Chapter 2)
- Similar to VoIP services (e.g., Vonage), but fully based on the individual peers
 - Skype requires a computer, VoIP services often do not
- Using resources: **Games, Video Streaming**; Controlling data: **OSN**;



P2P: Some Statistics

- Currently P2P accounts for 40% (*) of network traffic
 - A bit different in different networks
 - Hard to measure accurately
- Network providers (ISP) not too happy about this
 - But: Often traffic internal to ISP! (e.g., T-Com)
- Some numbers: (take with a grain of salt...)
 - KaZaA had 60 million users total, 1-5 million online at any time
 - 85 million downloads/day
 - Software downloaded over 230 million times
 - Google has 1 billion unique users / month, 400 million queries/day (50mio/200mio in 2006)
 - Skype has over 300 million users, over 20 million concurrently

(*) over 70% including file hosters and usenet

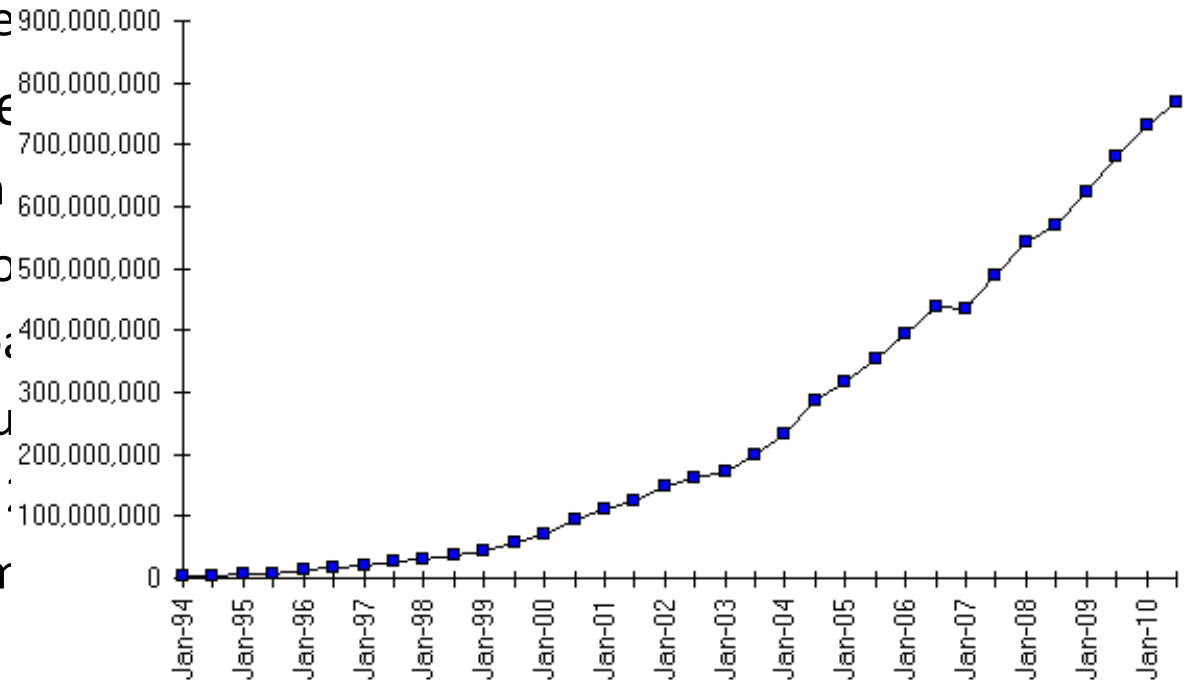


P2P: Some Statistics

- Currently P2P accounts for 40% (*) of network traffic
 - A bit different in different networks
 - Hard to measure accurately
- Network providers (ISP) not too happy about this

- But: Often traffic inte
- Some numbers: (take

- KaZaA had 60 million
 - 85 million downlo
 - Software downlo
- Google has 1 billion u
(50mio/200mio in
- Skype has over 300 m



(*) over 70% including file hosters and usenet



Why Does P2P Work?

- Why are P2P file sharing networks so successful?
 1. Easy to use
 - P2P software readily available, simple to use
 2. Provide something useful (for free)
 - Until recently, only alternative to P2P content was “buy a CD”
 - Online music stores may change this?
 3. Anyone can contribute
 - Contributions not tied to geographical location; user in Brazil can provide files for everyone (compare with ad hoc networks!)
 - Enough “altruistic” users to make P2P networks useful
- Some systems (Skype) completely hide the P2P-part
 - Will this become the future trend?



P2P: Traps and Pitfalls

- What could render current P2P networks useless?
 - In particular, file sharing networks
- 1. Removal of desirable content
 - Stricter enforcement of copyright laws?
- 2. Alternative ways of getting same content
 - Online music stores?
- 3. Blocking of P2P traffic by ISPs
 - Or making users pay for bandwidth they use?
- 4. Viruses or worms on P2P networks
 - Exploit bugs in P2P software
- 5. Frighten the users away...



When P2P and When Not P2P?

- So, when is P2P the right solution?
- Or, when is P2P the **wrong** solution?
- *Claim:* A general P2P vision is technically feasible
 - In other words, possible to build everything on Internet without any dedicated servers
- Gotcha: Just because it's technically feasible, it doesn't necessarily make sense...
- In other words, just because we can do it P2P, doesn't mean that we should do it P2P
 - True in many areas of life...
- So, when *is* P2P the right solution?!?



Some Criteria

- Let's consider the following criteria
 1. Budget
 - How much money do we have?
 2. Resource relevance
 - How widely are resources interesting to users?
 3. Trust
 - How much trust is there between users?
 4. Rate of system change
 - How fast does “something” in the system change
 5. Criticality
 - How critical is the service to the users



Analysis

Budget

- If you have enough money, build a centralized system
- Look at Google if you doubt this claim ;-)
 - Any system can be made to scale with enough money
- P2P is therefore useful when budget is not unlimited
 - In other words, most real-world situations...
 - From the rest of this analysis, we assume limited budget

Resource relevance

- If shared resources are highly relevant to a large number of users, P2P makes sense
- Easier to build a distributed solution when interest is widely spread



Analysis, Continued

Trust

- If other users can be trusted, P2P is a good solution
 - For example, corporate network or any closed network
- Building a fully distributed, **trusted** network is still very much a research problem (and may remain so...)

Rate of system change

Btw: what does „trust“ mean in this context!?

- How high are the system dynamics?
 - Rate of peers joining and leaving, rate of information change in system, rate of change in network topology, ...
- If the rate of change is too high, a distributed P2P solution might not be able to keep up
- Again, research problem



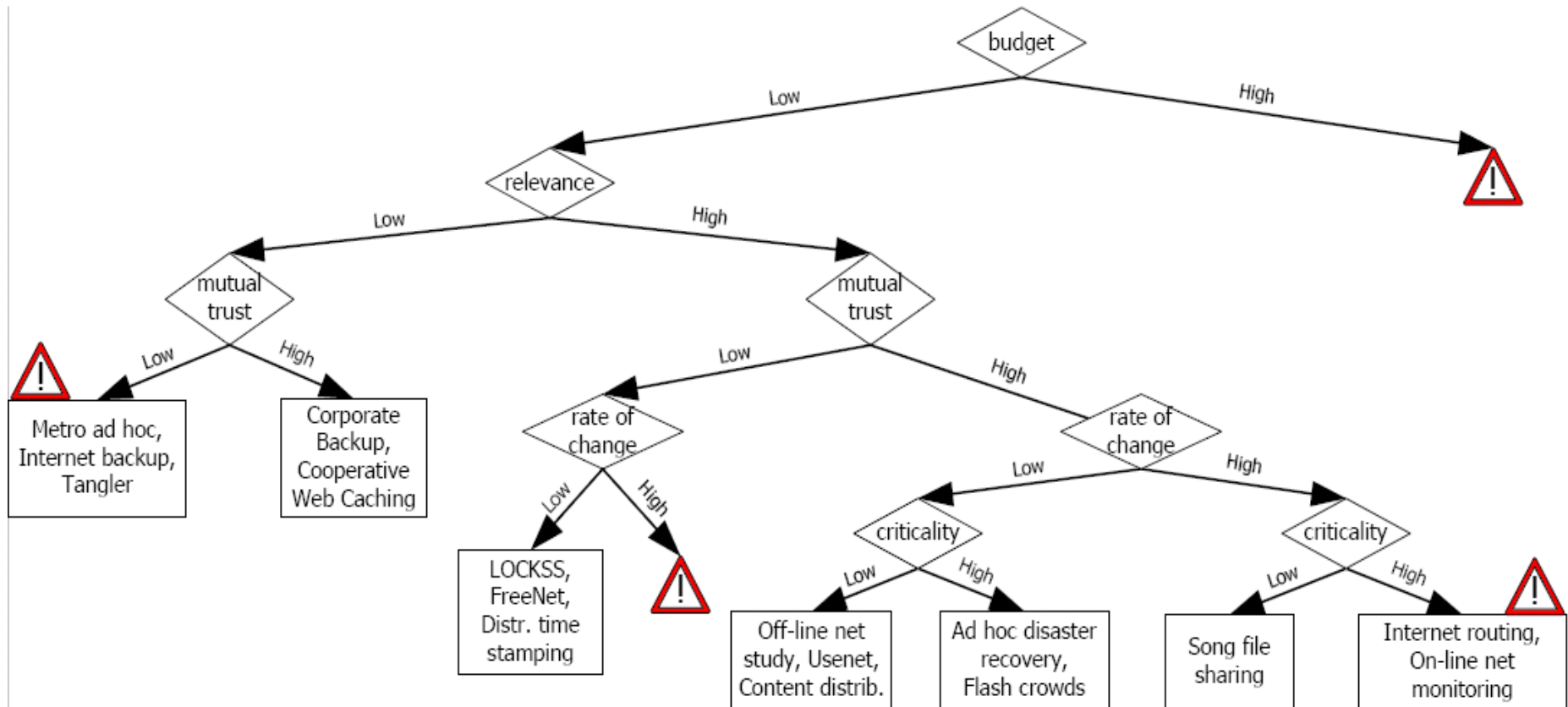
Analysis, End

Criticality

- How important is the service to the users?
- If you “can live without it”, P2P is acceptable
- If “it must work”, then consider other solutions...
- Summary: P2P is good when:
 - Budget is limited
 - Resources have wide interest and relevance
 - Trust between participants is high
 - Rate of change is manageable
 - Criticality is low
- Note: Again, no need to fulfill every point!



P2P Suitability Tree and Examples



Taken from M. Roussopoulos et al. "2 P2P or not 2 P2P?", IPTPS 2004



What does Future Hold for P2P?

- Take out crystal ball and look 5 years into future?
 - P2P has been around for just over 10 years now...
- Where will file sharing be in 5 years?
 - Still popular? Underground activity?
- P2P content distribution? (BitTorrent and others)
 - Microsoft building their system for software patches?
 - Some other systems patch via BitTorrent
- How about Skype and others?
 - Will Skype be around in 5 years?
 - Will Internet telephony be taken over by telcos?
- Research efforts in P2P?
 - More mature, concentrate on fundamental principles
 - What makes P2P different from other systems?



Future of P2P?

- Global P2P networks?
 - Besides file sharing, “Skype”, and research prototypes?
- Taking P2P concepts for other means and applications
 - Load balancing at S3 (inherently won)
 - Online Social Networking (remove central access to data)
 - Create resilient distributed systems (bot nets..)
- Insight on future trends: (at the example of Korea)
 - High bandwidth residential and wireless access
 - Online gaming (50% of network traffic!) main source of traffic
 - File sharing moved to pay models
 - Online communities gaining importance



Chapter Summary

- Peer-to-peer principle of self-organization and resource sharing
- Case Study of DNS to see it working the engineering way

- P2P systems exhibit specific characteristics:
 - Autonomy from central servers
 - Use of edge resources
 - Intermittent connectivity

- Hard to define clearly the limits of P2P
 - Quite some areas are closely related...
 - Different people working in different areas have different definitions



Outline of the Remainder of the Course

- Current P2P Systems
- Networks, Searching, and DHT
- Some Theory: Tools and Methods
- Novel Applications for P2P
 - Online Gaming
 - Online Social Networks
 - Application Level Multicast (P2P IPTV, Live Streaming)
 - P2P Botnets
- P2P and Security